# Data Cube: A Relational Aggregation Operator Generalizing Group-By, Cross-Tab, and Sub-Totals 

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How a relational database can support efficient extraction of multidimensional information

- What is a Data Cube in relational database?
- Why do we need the Cube Operator?
- How to implement the Cube Operator?
- What is a Data Cube in relational database?
- Why do we need the Cube Operator?
- How to implement the Cube Operator?


## Data Warehouse \& OLAP

Monitoring \& Admnistration


## Data Warehouse \& OLAP

## Multidimensional Data



## Data Cube

- A Multidimensional Data Model
- Dimension
- location
- time
- item
- Measurement
- sales

item (types)


## Data Cube

- In Relational Database:
- A relation with $n$-attribute domains

| Time | Item | Location | Sales |
| :--- | :--- | :--- | :--- |
| Q1 | Computer | Vancouver | 825 |
| Q1 | Security | Vancouver | 400 |
| Q2 | Phone | Vancouver | 31 |
| Q2 | Security | New York | 925 |
| Q3 | Security | Chicago | 789 |


item (types)

- What is a Data Cube in relational database?
- Why do we need the Cube Operator?
- How to implement the Cube Operator?


## Too Many Dimensions!

- Human are bad at understanding high dimensional data
- Need to reduce the dimension: super aggregation

item (type)

| time (quarter) | home <br> entertainment | computer | phone | security |
| :--- | :--- | :---: | :--- | :--- |
| Q1 | 605 | 825 | 14 | 400 |
| Q2 | 680 | 952 | 31 | 512 |
| Q3 | 812 | 1023 | 30 | 501 |
| Q4 | 927 | 1038 | 38 | 580 |

## Operations: Roll-Up

| Table 3.a: Sales Roll Up by Model by Year by Color |  |  |  |  |  |
| :--- | :---: | :---: | :---: | :---: | :---: |
| Model | Year | Color | Sales <br> by Model <br> by Year <br> by Color | Sales <br> by Model <br> by Year | Sales <br> by Model |
| Chevy | 1994 | black | 50 |  |  |
|  |  | white | 40 |  |  |
|  |  |  |  | 90 |  |
|  | 1995 | black | 85 |  |  |
|  |  | white | 115 |  |  |
|  |  |  |  | 200 |  |
|  |  |  |  |  | 290 |

Roll Up

## Operations: Roll-Up

| Table 3.a: Sales Roll Up by Model by Year by Color |  |  |  |  |  |
| :--- | :---: | :---: | :---: | :---: | :---: |
| Model | Year | Color | Sales <br> by Model <br> by Year <br> by Color | Sales <br> by Model <br> by Year | Sales <br> by Model |
| Chevy | 1994 | black | 50 |  |  |
|  |  | white | 40 |  |  |
|  |  |  |  | 90 |  |
|  | 1995 | black | 85 |  |  |
|  |  | white | 115 |  |  |
|  |  |  |  | 200 |  |
|  |  |  |  |  | 290 |


| Table 5.a: Sales Summary |  |  |  |
| :---: | :---: | :---: | :---: |
| Model | Year | Color | Units |
| Chevy | 1994 | black | 50 |
| Chevy | 1994 | white | 40 |
| Chevy | 1994 | ALL | 90 |
| Chevy | 1995 | black | 85 |
| Chevy | 1995 | white | 115 |
| Chevy | 1995 | ALL | 200 |
| Chevy | ALL | ALL | 290 |

`ALL` Value: Fill in the super-aggregation items

## Operation: Roll-Up

```
SELECT 'ALL', 'ALL', 'ALL', SUM(Sales)
    FROM Sales
    WHERE Model = 'Chevy'
UNION
SELECT Model, 'ALL', 'ALL', SUM(Sales)
    FROM Sales
    WHERE Model = 'Chevy'
    GROUP BY Model
UNION
SELECT Model, Year, 'ALL', SUM(Sales)
    FROM Sales
    WHERE Model = 'Chevy'
    GROUP BY Model, Year
UNION
SELECT Model, Year, Color, SUM(Sales)
    FROM Sales
    WHERE Model = 'Chevy'
    GROUP BY Model, Year, Color;
```

| Table 5.a: Sales Summary |  |  |  |
| :---: | :---: | :---: | :---: |
| Model | Year | Color | Units |
| Chevy | 1994 | black | 50 |
| Chevy | 1994 | white | 40 |
| Chevy | 1994 | ALL | 90 |
| Chevy | 1995 | black | 85 |
| Chevy | 1995 | white | 115 |
| Chevy | 1995 | ALL | 200 |
| Chevy | ALL | ALL | 290 |
| ALL |  |  |  |

N dimensions: N Unions

## Operation: Roll-Up

## Asymmetric! <br> Missing:

| Table 5.b: Sales Summary rows missing form |  |  |  |
| :---: | :---: | :---: | :---: |
| Table 5.a to convert the roll-up into a cube. |  |  |  |
| Model | Year | Color | Units |
| Chevy | ALL | black | 135 |
| Chevy | ALL | white | 155 |


| Table 5.a: Sales Summary |  |  |  |
| :---: | :---: | :---: | :---: |
| Model | Year | Color | Units |
| Chevy | 1994 | black | 50 |
| Chevy | 1994 | white | 40 |
| Chevy | 1994 | ALL | 90 |
| Chevy | 1995 | black | 85 |
| Chevy | 1995 | white | 115 |
| Chevy | 1995 | ALL | 200 |
| Chevy | ALL | ALL | 290 |
| ALL |  |  |  |

```
UNION
SELECT Model, 'ALL', Color, SUM(Sales)
    FROM Sales
    WHERE Model = 'Chevy'
    GROUP BY Model, Color;
```


## Operation: Cross Tab

- Cross Tabulation: 2D Symmetric Aggregation Result

| Table 6.a: Chevy Sales Cross Tab |  |  |  |
| :---: | :---: | :---: | :---: |
| Chevy | $\mathbf{1 9 9 4}$ | $\mathbf{1 9 9 5}$ | total (ALL) |
| black | 50 | 85 | 135 |
| white | 40 | 115 | 155 |
| total (ALL) | 90 | 200 | 290 |

## 3-D Generalization of Cross Tab



## 3-D Generalization of Cross Tab

## Aggregate



```
SELECT 'ALL', `ALL', 'ALL', SUM(Sales)
UNION
SELECT Model, 'ALL', `ALL', SUM(Sales)
    FROM Sales
GROUP BY Model
```


## Group By x 8

```
UNION
SELECT Model, Year, 'ALL', SUM(Sales)
    FROM Sales
    GROUP BY Model, Year
UNION
SELECT Model, Year, Color, SUM(Sales)
    FROM Sales
GROUP BY Model, Year, Color;
```


## N-D Generalization of Cross Tab

Aggregate

SELECT `ALL', `ALL', 'ALL', SUM(Sales)
SELECT `ALL', `ALL', 'ALL', SUM(Sales)
UNION
UNION
SELECT Model, 'ALL', 'ALL', SUM(Sales)
SELECT Model, 'ALL', 'ALL', SUM(Sales)
FROM Sales
FROM Sales
GROUP BY Model
GROUP BY Model

## Group By $\times 2^{N}$

UNION
UNION
SELECT Model, Year, 'ALL', SUM(Sales)
SELECT Model, Year, 'ALL', SUM(Sales)
FROM Sales
FROM Sales
GROUP BY Model, Year
GROUP BY Model, Year
UNION
UNION
SELECT Model, Year, Color, SUM(Sales)
SELECT Model, Year, Color, SUM(Sales)
FROM Sales
FROM Sales
GROUP BY Model, Year, Color;
GROUP BY Model, Year, Color;

## N-D Generalization of Cross Tab

- Problems
- Expressing with conventional SQL is exhaustive
- Too complex to analyze for optimization


## The Cube Operator



- What is a Data Cube in relational database?
- Why do we need the Cube Operator?
- How to implement the Cube Operator?
- How the Cube fit in SQL?
- How to compute the Cube?


## Compute The Cube

```
SELECT `ALL', `ALL', 'ALL', SUM(Sales)
```

SELECT `ALL', `ALL', 'ALL', SUM(Sales)

```
SELECT `ALL', `ALL', 'ALL', SUM(Sales)
    FROM Sales
    FROM Sales
    FROM Sales
UNION
UNION
UNION
SELECT Model, 'ALL', 'ALL', SUM(Sales)
SELECT Model, 'ALL', 'ALL', SUM(Sales)
SELECT Model, 'ALL', 'ALL', SUM(Sales)
    FROM Sales
    FROM Sales
    FROM Sales
GROUP BY Model
```

GROUP BY Model

```
GROUP BY Model
```

```
SELECT Model, Year, Color, SUM(sales) AS Sales
FROM Sales
WHERE Model in {'Ford', 'Chevy'}
    AND Year BETWEEN 1990 AND 1992
GROUP BY CUBE Model, Year, Color;
=
```

UNION
SELECT Model, Year, 'ALL', SUM(Sales)
FROM Sales
GROUP BY Model, Year
UNION
SELECT Model, Year, Color, SUM(Sales)
FROM Sales
GROUP BY Model, Year, Color;

## Implementation of Aggregate Functions

- Start = Init (\&handle):Allocates the handle and initializes the aggregate computation
- Next - Iter (\&handle, value): Aggregates the next value into the current aggregate
- End - Final(\&handle): Computes and returns the resulting aggregate by using data saved in the handle. This invocation deallocates the handle


## Implementation of Aggregate Functions

## Example of AVG



- Start - Init (\&handle):Allocates the handle and initializes the aggregate computation
- Next - Iter (\&handle, value): Aggregates the next value into the current aggregate
- End - Final(\&handle):

Computes and returns the resulting aggregate by using data saved in the handle. This invocation deallocates the handle

## Implementation of Aggregate Functions

## Example of AVG

```
(..., value)
next handle
sum+=value count+=1
```

- Start - Init (\&handle):Allocates the handle and initializes the aggregate computation
- Next - Iter (\&handle, value):

Aggregates the next value into the current aggregate

- End - Final(\&handle):

Computes and returns the resulting aggregate by using data saved in the handle. This invocation deallocates the handle

## Implementation of Aggregate Functions

## Example of AVG

handle
count

sum/count

- Start - Init (\&handle):Allocates the handle and initializes the aggregate computation
- Next - Iter (\&handle, value): Aggregates the next value into the current aggregate
- End - Final(\&handle): Computes and returns the resulting aggregate by using data saved in the handle. This invocation deallocates the handle


## $2^{N}$ Algorithm

I. Allocate a handle for each cell of the cube - Init()
2. Each tuple needs to invoke the Iter() function once for the cells that match the tuple
3. Compute result for each cell of the cube - Final()

| Branch | Model | Year | Color | Sales |
| :--- | :--- | :--- | :--- | :--- |
| Burnaby | Chevy | 1990 | red | 23 |
| Richmond | Chevy | 1990 | white | 14 |
| Richmond | Chevy | 1990 | white | 31 |
| Burnaby | Ford | 1990 | blue | 23 |
| Richmond | Ford | 1990 | red | 4 |
| Burnaby | Chevy | 1991 | blue | 22 |
| Richmond | Ford | 1992 | red | 32 |
| $\ldots$ | $\ldots$ | $\ldots$ | $\ldots$ | $\ldots$ |



## $2^{N}$ Algorithm

- Invoke Init() \& Final() one time for each cell
- Invoke Iter() $2^{N}$ times for each tuple

| Branch | Model | Year | Color | Sales |
| :--- | :--- | :--- | :--- | :--- |
| Burnaby | Chevy | 1990 | red | 23 |
| Richmond | Chevy | 1990 | white | 14 |
| Richmond | Chevy | 1990 | white | 31 |
| Burnaby | Ford | 1990 | blue | 23 |
| Richmond | Ford | 1990 | red | 4 |
| Burnaby | Chevy | 1991 | blue | 22 |
| Richmond | Ford | 1992 | red | 32 |
| $\ldots$ | $\ldots$ | $\ldots$ | $\ldots$ | $\ldots$ |



## $2^{N}$ Algorithm

- Invoke Init() \& Final() one time for each cell
- Invoke Iter() $2^{N}$ times for each tuple
$\longleftarrow$ Can be optimized

| Branch | Model | Year | Color | Sales |
| :--- | :--- | :--- | :--- | :--- |
| Burnaby | Chevy | 1990 | red | 23 |
| Richmond | Chevy | 1990 | white | 14 |
| Richmond | Chevy | 1990 | white | 31 |
| Burnaby | Ford | 1990 | blue | 23 |
| Richmond | Ford | 1990 | red | 4 |
| Burnaby | Chevy | 1991 | blue | 22 |
| Richmond | Ford | 1992 | red | 32 |
| $\ldots$ | $\ldots$ | $\ldots$ | $\ldots$ | $\ldots$ |



## Computing the Cube

- Speed Up the Process
- Make use of the middle result:

> N-D Aggregate -> (N-I)-D Aggregate

| Branch | Model | Year | Color | Sales |
| :--- | :--- | :--- | :--- | :--- |
| Burnaby | Chevy | 1990 | red | 23 |
| Richmond | Chevy | 1990 | white | 14 |
| Richmond | Chevy | 1990 | white | 31 |
| Burnaby | Ford | 1990 | blue | 23 |
| Richmond | Ford | 1990 | red | 4 |
| Burnaby | Chevy | 1991 | blue | 22 |
| Richmond | Ford | 1992 | red | 32 |
| $\ldots$ | $\ldots$ | $\ldots$ | $\ldots$ | $\ldots$ |



## Computing the Cube

## Can F be computed in distributive manner?



| Branch | Model | Year | Color | Sales |
| :--- | :--- | :--- | :--- | :--- |
| Burnaby | Chevy | 1990 | red | 23 |
| Richmond | Chevy | 1990 | white | 14 |
| Richmond | Chevy | 1990 | white | 31 |
| Burnaby | Ford | 1990 | blue | 23 |
| Richmond | Ford | 1990 | red | 4 |
| Burnaby | Chevy | 1991 | blue | 22 |
| Richmond | Ford | 1992 | red | 32 |
| $\ldots$ | $\ldots$ | $\ldots$ | $\ldots$ | $\ldots$ |



## Aggregate Functions Classification

- Distributive: SUM(), MIN(), MAX(), COUNT()
- Can be computed in a distributive manner
- Algebraic:AVG(), $\operatorname{MaxN}(), \operatorname{MinN}()$
- Can be computed in a distributive manner with $m$ arguments - need to keep both the handle \& the result for each cell
- Holistic: Median()
- No constant m exists - need to scan all the tuples


## Computing the Cube

- Speed Up the Process for Distributive \& Algebraic Functions
- Make use of the middle result

Aggregate on the smallest list


## Maintaining The Cube

- Trigger Conditions: UPDATE, INSERT, DELETE
- Can be different for the same function: MAX()
- INSERT: Distributive
- DELETE/UPDATE: Holistic

| Branch | Model | Year | Color | Sales |
| :--- | :--- | :--- | :--- | :--- |
| Burnaby | Chevy | 1990 | red | 23 |
| Richmon | Chevy | 1990 | white | 14 |
| Richmon | Chevy | 1990 | white | 31 |
| Burnaby | Ford | 1990 | blue | 23 |
| Richmon | Ford | 1990 | red | 4 |
| Burnaby | Chevy | 1991 | blue | 22 |
| Richmon | Ford | 1992 | red | 32 |
| $\ldots$ | $\ldots$ | $\ldots$ | $\ldots$ | $\ldots$ |



## TakeAways

- The cube operator computes aggregations over all possible subsets of the specified dimensions
- The result of the cube operator can be modeled as a data cube
- We can speed up the computation of the cube for many common aggregate functions by using the middle result

Q \& A

